

## Spectral analysis of Wikipedia and PhysRev networks

## Klaus Frahm

Quantware MIPS Center
Université Paul Sabatier
Laboratoire de Physique Théorique, UMR 5152, IRSAMC, CNRS
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## Google matrix for directed

## networks

Define the adjacency matrix $A$ by $A_{i j}=1$ if there is a link from the node $j$ to $i$ in the network (of size $N$ ) and $A_{i j}=0$ otherwise. Let $S_{i j}=A_{i j} / \sum_{i} A_{i j}$ and $S_{i j}=1 / N$ if $\sum_{i} A_{i j}=0$ (dangling nodes). $S$ is of Perron-Frobenius type but for many networks the eigenvalue $\lambda_{1}=1$ is highly degenerate [ $\Rightarrow$ convergence problem to arrive at the stationary limit of $p(t+1)=S p(t)$ ].
Therefore define the Google matrix:

$$
G(\alpha)=\alpha S+(1-\alpha) \frac{1}{N} e e^{T}
$$

where $e=(1, \ldots, 1)^{T}$ and $\alpha=0.85$ is a typical damping factor. Here there is a unique eigenvector for $\lambda_{1}=1$ called the PageRank $P$ and the convergence goes with $\alpha^{t}$.
(CheiRank $P^{*}$ by replacing: $A \rightarrow A^{*}=A^{T}$ ).

## Arnoldi method

to (partly) diagonalize large sparse non-symmetric $d \times d$ matrices:

- choose an initial normalized vector $\xi_{0}$ (random or "otherwise")
- determine the Krylov space of dimension $n_{A}$ (typically:
$\left.1 \ll n_{A} \ll d\right)$ spanned by the vectors: $\xi_{0}, G \xi_{0}, \ldots, G^{n_{A}-1} \xi_{0}$
- determine by Gram-Schmidt orthogonalization an orthonormal basis $\left\{\xi_{0}, \ldots, \xi_{n_{A}-1}\right\}$ and the representation of $G$ in this basis:

$$
G \xi_{k}=\sum_{j=0}^{k+1} H_{j k} \xi_{j}
$$

- diagonalize the Arnoldi matrix $H$ which has Hessenberg form:
$H=\left(\begin{array}{ccccc}* & * & \cdots & * & * \\ * & * & \cdots & * & * \\ 0 & * & \cdots & * & * \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & * & * \\ \hline 0 & 0 & \cdots & 0 & *\end{array}\right)$ which provides the Ritz eigenvalues that are
very good aproximations to the "largest" eigenvalues of $A$.


## Invariant subspaces

In realistic WWW or other networks invariant subspaces of nodes create (possibly) large degeneracies of $\lambda_{1}$ (or $\lambda_{2}$ if $\alpha<1$ ) which is very problematic for the Arnoldi method.
Therefore one needs to determine the invariant subspaces defined as subsets of nodes such that for any node in a subspace each outgoing link stays in the subspace. One can efficiently find all subspaces of maximal size (or dimension) $N_{c}$ (with $N_{c}=b N$ a certain fraction of the network size $N$, e.g. $b=0.1$ ) and then all subspaces with common members are merged resulting in a decomposition of the network in many separate subspaces with $N_{s}$ nodes and a "big" core space of the remaining $N-N_{s}$ nodes.
Note that dangling nodes are by construction core space nodes.
Possible: core space node $\rightarrow$ subspace node
Impossible: subspace node $\rightarrow$ core space node

The decomposition in subspaces and a core space implies a block structure of the matrix $S$ :

$$
S=\left(\begin{array}{cc}
S_{s s} & S_{s c} \\
0 & S_{c c}
\end{array}\right) \quad, \quad S_{s s}=\left(\begin{array}{ccc}
S_{1} & 0 & \cdots \\
0 & S_{2} & \\
\vdots & & \ddots
\end{array}\right)
$$

where $S_{s s}$ is block diagonal according to the subspaces. The subspace blocks of $S_{s s}$ are all matrices of PF type with at least one eigenvalue $\lambda_{1}=1$ explaining the high degeneracies.
To determine the spectrum of $S$ apply:

- Exact (or Arnoldi) diagonalization on each subspace.
- The Arnoldi method to $S_{c c}$ to determine the largest core space eigenvalues $\lambda_{j}$ (note: $\left|\lambda_{j}\right|<1$ ). The largest eigenvalues of $S_{c c}$ are no longer degenerate but other degeneracies are possible (e.g. $\lambda_{j}=0.9$ for Wikipedia).


## Spectrum of Wikipedia

L. Ermann, KMF and D.L. Shepelyansky, Eur. Phys. J. B 86, 193 (2013) Wikipedia 2009 : $N=3282257$ nodes, $N_{\ell}=71012307$ network links.


$n_{A}=6000$ for both cases

## Some Eigenvectors:









left (right): PageRank (CheiRank)
black: PageRank (CheiRank) at $\alpha=0.85$
grey: PageRank (CheiRank) at $\alpha=1-10^{-8}$
red and green: first two core space eigenvectors
blue and pink: two eigenvectors with large imaginary part in the eigenvalue

## Detail study of 200 selected eigenvectors

 with eigenvalues "close" to the unit circle:

## Power law decay of eigenvectors:






$$
\begin{gathered}
\left|\psi_{i}\left(K_{i}\right)\right| \sim K_{i}^{b} \quad \text { for } \quad K_{i} \geq 10^{4} \\
\varphi=\arg \left(\lambda_{i}\right)
\end{gathered}
$$

## Inverse participation ratio of eigenvectors:






$$
\begin{gathered}
\xi_{\mathrm{IPR}}=\left(\sum_{j}\left|\psi_{i}(j)\right|^{2}\right)^{2} / \sum_{j}\left|\psi_{i}(j)\right|^{4} \\
\varphi=\arg \left(\lambda_{i}\right)
\end{gathered}
$$

## "Themes" of certain eigenvectors:



Number of links between or inside sets $A$ and $B$ defined by the index $K_{i}$ ordered by decreasing absolute value of Wikipedia eigenstates:


## Physical Review network

(work in progress: KMF, Young-Ho Eom, D. Shepelyansky)
$N=463347$ nodes and $N_{\ell}=4691015$ links.
Coarse-grained matrix structure ( $500 \times 500$ cells):

left: time ordered
right: journal and then time ordered
"11" Journals of Physical Review: (Phys. Rev. Series I), Phys. Rev., Phys. Rev. Lett., (Rev. Mod. Phys.), Phys. Rev. A, B, C, D, E, (Phys. Rev. STAB and Phys. Rev. STPER).
$\Rightarrow$ nearly triangular matrix structure of adjancy matrix: most citations links $t \rightarrow t^{\prime}$ are for $t>t^{\prime}$ ("past citations") but there is small number (12126 $=2.6 \times 10^{-3} N_{\ell}$ ) of links $t \rightarrow t^{\prime}$ with $t \leq t^{\prime}$ corresponding to future citations.
Spectrum by "double-precision" Arnoldi method with $n_{A}=8000$ :


Numerical problem: eigenvalues with $|\lambda|<0.3-0.4$ are not reliable! Reason: large Jordan subspaces associated to the eigenvalue $\lambda=0$.
"very bad" Jordan perturbation theory:
Consider a "perturbed" Jordan block of size $D$ :

$$
\left(\begin{array}{ccccc}
0 & 1 & \cdots & 0 & 0 \\
0 & 0 & \cdots & 0 & 0 \\
\vdots & \vdots & \ddots & \vdots & \vdots \\
0 & 0 & \cdots & 0 & 1 \\
\varepsilon & 0 & \cdots & 0 & 0
\end{array}\right)
$$

characteristic polynomial: $\lambda^{D}-(-1)^{D} \varepsilon$
$\varepsilon=0 \quad \Rightarrow \quad \lambda=0$
$\varepsilon \neq 0 \quad \Rightarrow \quad \lambda_{j}=-\varepsilon^{1 / D} \exp (2 \pi i j / D)$
for $D \approx 10^{2}$ and $\varepsilon=10^{-16} \Rightarrow$ "Jordan-cloud" of artifical eigenvalues due to rounding errors in the region $|\lambda|<0.3-0.4$.

## Triangular approximation

Remove the small number of links due to "future citations".
Semi-analytical diagonalization is possible:

$$
S=S_{0}+e d^{T} / N
$$

where $e_{n}=1$ for all nodes $n, d_{n}=1$ for dangling nodes $n$ and $d_{n}=0$ otherwise. $S_{0}$ is the pure link matrix which is nil-potent:

$$
S_{0}^{l}=0 \text { with } l=352
$$

Let $\psi$ be an eigenvector of $S$ with eigenvalue $\lambda$ and $C=d^{T} \psi$.

- If $C=0 \Rightarrow \psi$ eigenvector of $S_{0} \Rightarrow \lambda=0$ since $S_{0}$ nil-potent. These eigenvectors belong to large Jordan blocks and are responsible for the numerical problems.
Note: Similar situation as in network of integer numbers where $l=\left[\log _{2}(N)\right]$ and numerical instability for $|\lambda|<0.01$.
- If $C \neq 0 \Rightarrow \lambda \neq 0$ since the equation $S_{0} \psi=-C e / N$ does not have a solution $\Rightarrow \lambda 1-S_{0}$ invertible.

$$
\begin{gathered}
\Rightarrow \psi=C\left(\lambda \mathbf{1}-S_{0}\right)^{-1} e / N=\frac{C}{\lambda} \sum_{j=0}^{l-1}\left(\frac{S_{0}}{\lambda}\right)^{j} e / N . \\
\text { From } \lambda^{l}=\left(d^{T} \psi / C\right) \lambda^{l} \Rightarrow \mathcal{P}_{r}(\lambda)=0
\end{gathered}
$$

with the reduced polynomial of degree $l=352$ :

$$
\mathcal{P}_{r}(\lambda)=\lambda^{l}-\sum_{j=0}^{l-1} \lambda^{l-1-j} c_{j}=0 \quad, \quad c_{j}=d^{T} S_{0}^{j} e / N
$$

$\Rightarrow$ at most $l=352$ eigenvalues $\lambda \neq 0$ which can be numerically determined as the zeros of $\mathcal{P}_{r}(\lambda)$.
However: still numerical problems:

- $c_{l-1} \approx 3.6 \times 10^{-352}$
- alternate sign problem with a strong loss of significance.
- big sensitivity of eigenvalues on $c_{j}$


## Solution:

Using the multi precision library GMP with 256 binary digits the zeros of $\mathcal{P}_{r}(\lambda)$ can be determined with accuracy $\sim$ $10^{-18}$.
Furthermore the Arnoldi method can also be implemented with higher precision.
red crosses: zeros of $\mathcal{P}_{r}(\lambda)$ from 256 binary digits calculation
blue squares: eigenvalues from Arnoldi method with $52,256,512,1280$ binary digits. In the last case: $\Rightarrow$ break off at $n_{A}=352$ with vanishing coupling element.


## Full Physical Review network

High precision Arnoldi method for full Physical Review network (including the "future citations") for 52, 256, 512, 768 binary digits and $n_{A}=2000$ :


## Degeneracies



High precision in Arnoldi method is "bad" to count the degeneracy of certain degenerate eigenvalues.

In theory the Arnoldi method cannot find several eigenvectors for degenerate eigenvalues, a shortcoming which is (partly) "repaired" by rounding errors.

Q: How are highly degenerate core space eigenvalues possible?

## Semi-analytical argument for the full PR network:

$$
S=S_{0}+e d^{T} / N
$$

There are two groups of eigenvectors $\psi$ with: $S \psi=\lambda \psi$

1. Those with $d^{T} \psi=0 \Rightarrow \psi$ is also an eigenvector of $S_{0}$. Generically an arbitrary eigenvector of $S_{0}$ is not an eigenvector of $S$ unless the eigenvalue is degenerate with degeneracy $m>1$. Using linear combinations of different eigenvectors for the same eigenvalue one can construct $m-1$ eigenvectors $\psi$ respecting $d^{T} \psi=0$ which are therefore eigenvectors of $S$.
Pratically: determine degenerate subspace eigenvalues of $S_{0}$ (and also of $S_{0}^{T}$ ) which are of the form: $\lambda= \pm 1 / \sqrt{n}$ with $n=1,2,3, \ldots$ due to $2 \times 2$-blocks:

$$
\left(\begin{array}{cc}
0 & 1 / n_{1} \\
1 / n_{2} & 0
\end{array}\right) \Rightarrow \lambda= \pm \frac{1}{\sqrt{n_{1} n_{2}}} .
$$

2. Those with $d^{T} \psi \neq 0 \quad \Rightarrow \quad \mathcal{R}(\lambda)=0$ with the rational function:

$$
\mathcal{R}(\lambda)=1-d^{T} \frac{1}{\lambda \mathbf{1}-S_{0}} e / N=1-\sum_{j, q} \frac{C_{j q}}{\left(\lambda-\rho_{j}\right)^{q}}
$$

Here $C_{j q}$ and $\rho_{j}$ are unknown, except for
$\rho_{1}=2 \operatorname{Re}\left[(9+i \sqrt{119})^{1 / 3}\right] /(135)^{1 / 3} \approx 0.9024$ and
$\rho_{2,3}= \pm 1 / \sqrt{2} \approx \pm 0.7071$.
Idea: Expand the geometric matrix series $\Rightarrow$

$$
\mathcal{R}(\lambda)=1-\sum_{j=0}^{\infty} c_{j} \lambda^{-1-j} \quad, \quad c_{j}=d^{T} S_{0}^{j} e / N
$$

which converges for $|\lambda|>\rho_{1} \approx 0.9024$ since $c_{j} \sim \rho_{1}^{j}$ for $j \rightarrow \infty$.

Problem: How to determine the zeros of $\mathcal{R}(\lambda)$ with $|\lambda|<\rho_{1}$ ?

## Analytic continuation by rational interpolation:

Use the series to evaluate $\mathcal{R}(z)$ at $n_{S}$ support points
$z_{j}=\exp \left(2 \pi i j / n_{S}\right)$ with a given precision of $p$ binary digits and determine the rational function $R_{I}(z)$ which interpolates $\mathcal{R}(z)$ at these support points. Two cases:

$$
\begin{aligned}
& n_{S}=2 n_{R}+1 \quad \Rightarrow \quad R_{I}(z)=\frac{P_{n_{R}}(z)}{Q_{n_{R}}(z)} \\
& n_{S}=2 n_{R}+2 \quad \Rightarrow \quad R_{I}(z)=\frac{P_{n_{R}}(z)}{Q_{n_{R}+1}(z)}
\end{aligned}
$$

The $n_{R}$ zeros of $P_{n_{R}}(z)$ are approximations of the eigenvalues of $S$ (of the 2nd group).
For a given precision, e. g. $p=1024$ binary digits one can obtain a certain number of reliable eigenvalues, e. g. $n_{R}=300$. The method can be pushed up to $p=16384$ and $n_{R}=2500$ which is better than the high precision Arnoldi method with $n_{A}=2000$.

## Examples:

Some "artificial zeros" for $n_{R}=340$ and $p=1024$ (left top and middle pane/s) where both variants of the method differ.

For $n_{R}=300$ and $p=1024$ most zeros coincide with HP Arnoldi method (right top and middle panels) and both variants of the method coincide.

Lower panels: comparison for $n_{R}=$ 2000, $p=12288$ (left) or for $n_{R}=$ $2500, p=16384$ with HP Arnoldi method.


Accurate eigenvalue spectrum for the full Physical Review network by the rational interpolation method (left) and the HP Arnoldi method (right):




## Conclusion

- Detailed eigenvector study for the Wikipedia network.
- Identification of certain themes or communities with the help of eigenvectors.
- Subtle numerical problems for the eigenvalue problem of the Physical Review network which can be solved by a semi-analytical method and a high precision implementation of the Arnoldi method.
- Understanding of the degeneracies of core space eigenvalues and a decompostion of the core space eigenvalues in two groups. Important role of subspaces of $S_{0}$ (very different from the subspaces of $S$ !).
- New rational interpolation method to determine accurately the eigenvalues of a network matrix. Well suited for nearly triangular matrices but works in principle also for other case (e. g. Wikipedia but less efficient here).
- Drastic effect of the triangular approximation on the eigenvalue spectrum. Strong reduction of non-vanishing eigenvalues, from about $\sim 8000-10000$ to 352 and only very few eigenvalues on the real axis. This implies a very strong effect of the few future citations on the spectrum.
- Very useful applications of the GNU high precision library GMP: http://gmplib.org/ for different numerical methods: determination of zeros of the reduced polynomial, rational interpolation method, Arnoldi method.


## Appendix:

The subspace of $\lambda \neq 0$ is represented by the vectors $v^{(j)}=S_{0}^{j-1} e / N$ for $j=1, \ldots, l$

$$
\Rightarrow \quad S v^{(j)}=c_{j-1} v^{(1)}+v^{(j+1)}=\sum_{k=0}^{l-1} \bar{S}_{k, j} v^{(k)}
$$

"Small" $l \times l$-representation matrix :

$$
\bar{S}=\left(\begin{array}{ccccc}
c_{0} & c_{1} & \cdots & c_{l-2} & c_{l-1} \\
1 & 0 & \cdots & 0 & 0 \\
0 & 1 & \cdots & 0 & 0 \\
\vdots & \vdots & \ddots & \vdots & \vdots \\
0 & 0 & \cdots & 1 & 0
\end{array}\right) \quad, \quad \bar{P}=C\left(\begin{array}{c}
1 \\
1 \\
1 \\
\vdots \\
1
\end{array}\right)
$$

with $P=\sum_{j} \bar{P}_{j} v^{(j)}=C \sum_{j} v^{(j)}$ and due to sum rule: $\sum_{j} c_{j}=1$.

