



#### **Google matrix of directed networks**

#### **Klaus Frahm**

#### **Quantware MIPS Center**

Université Paul Sabatier Laboratoire de Physique Théorique, UMR 5152, IRSAMC, CNRS supported by EC FET Open project NADINE

Trento Workshop, Spectral properties of complex networks, Trento, 24 July 2012

## **Perron-Frobenius operators**

discrete Markov process:

$$p_i(t+1) = \sum_j G_{ij} p_j(t)$$

with probabilities  $p_i(t) \ge 0$  and the *Perron-Frobenius* matrix G such that:

$$\sum_{i} G_{ij} = 1 \quad , \quad G_{ij} \ge 0 \; .$$

For any vector v :

$$\Rightarrow \qquad \|Gv\|_1 \le \|v\|_1$$

 $\Rightarrow$  complex eigenvalues  $|\lambda_j| \le 1$  and (at least) one eigenvalue  $\lambda_1 = 1$  and its right eigenvector P is the stationary distribution:

$$P = \lim_{t \to \infty} p(t)$$

provided  $\lambda_1$  is not degenerate !

## Google matrix for directed networks

Define the *adjacency matrix* A by  $A_{ij} = 1$  if there is a link from the node j to i in the network (of size N) and  $A_{ij} = 0$  otherwise. Let  $S_{ij} = A_{ij} / \sum_i A_{ij}$  and  $S_{ij} = 1/N$  if  $\sum_i A_{ij} = 0$  (dangling nodes). S is of Perron-Frobenius type but for many networks the eigenvalue  $\lambda_1 = 1$  is highly degenerate [ $\Rightarrow$  convergence problem to arrive at the stationary limit of p(t+1) = S p(t)].

Therefore define the **Google matrix**:

$$G(\alpha) = \alpha S + (1 - \alpha) \frac{1}{N} e e^{T}$$

where  $e = (1, ..., 1)^T$  and  $\alpha = 0.85$  is a typical damping factor. Here there is unique eigenvector for  $\lambda_1 = 1$  called the *PageRank* P and the convergence goes with  $\alpha^t$ .

(**CheiRank**  $P^*$  by replacing:  $A \to A^* = A^T$ ).

## **Ulam Method**

(Ermann, Shepelyansky (2010), KF, Shepelyansky (2010))

to construct a *Perron-Frobenius matrix* as discrete approximation for the PF operator of dynamical systems with mixed phase space:

- Subdivide phase space in discrete cells.
- Iterate (for a very long time) a classical trajectory and attribute a new number to each new cell which is entered for the first time. At the same time count the number of transitions from cell *i* to cell *j* (⇒ n<sub>ji</sub>).

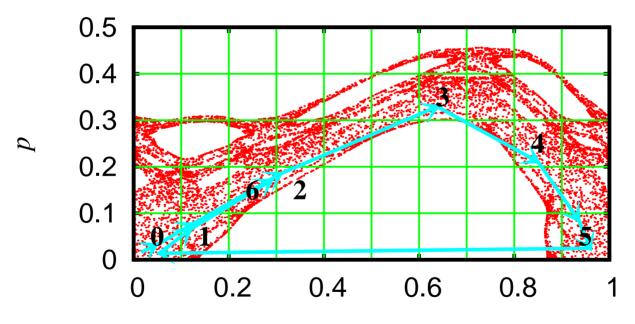
 $\bullet \Rightarrow$  The matrix

$$G_{ji} = \frac{n_{ji}}{\sum_l n_{li}}$$

is of Perron-Frobenius type :  $G_{ji} \ge 0$ ,  $\sum_j G_{ji} = 1$ .

#### **Chirikov Standard map**

$$p_{n+1} = p_n + \frac{k}{2\pi} \sin(2\pi x_n)$$
  
$$x_{n+1} = x_n + p_{n+1} , \quad k = k_c = 0.971635406$$



#### Arnoldi method

to (partly) diagonalize large sparse non-symmetric  $d \times d$  matrices:

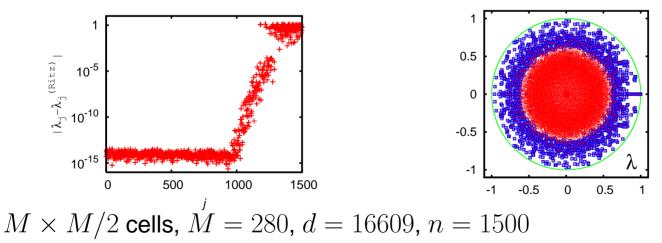
- choose an initial normalized vector  $\xi_0$  (random or "otherwise")
- determine the *Krylov space* of dimension n (typically:  $1 \ll n \ll d$ ) spanned by the vectors:  $\xi_0, G \xi_0, \ldots, G^{n-1} \xi_0$
- determine by *Gram-Schmidt* orthogonalization an orthonormal basis  $\{\xi_0, \ldots, \xi_{n-1}\}$  and the representation of *G* in this basis:

$$G\,\xi_k = \sum_{j=0}^{k+1} H_{jk}\,\xi_j$$

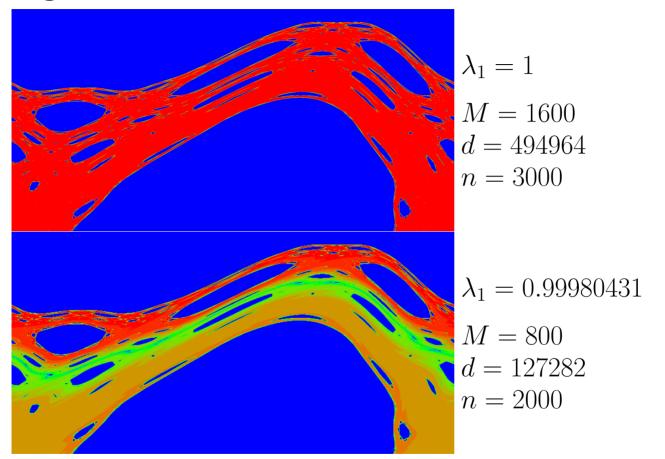
• diagonalize the **Arnoldi matrix** H which has **Hessenberg** form:

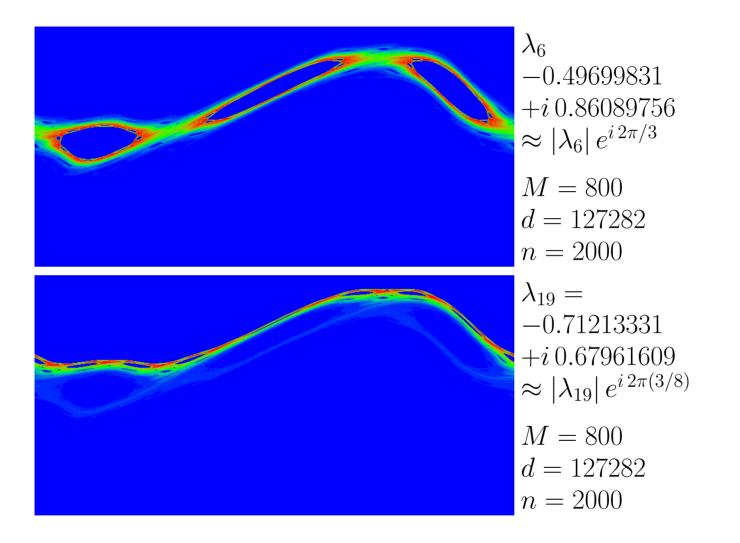
$$H = \begin{pmatrix} * & * & \cdots & * & * \\ * & * & \cdots & * & * \\ 0 & * & \cdots & * & * \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & * & * \\ \hline 0 & 0 & \cdots & 0 & * \end{pmatrix}$$

which provides the *Ritz eigenvalues* that are very good aproximations to the "largest" eigenvalues of A.



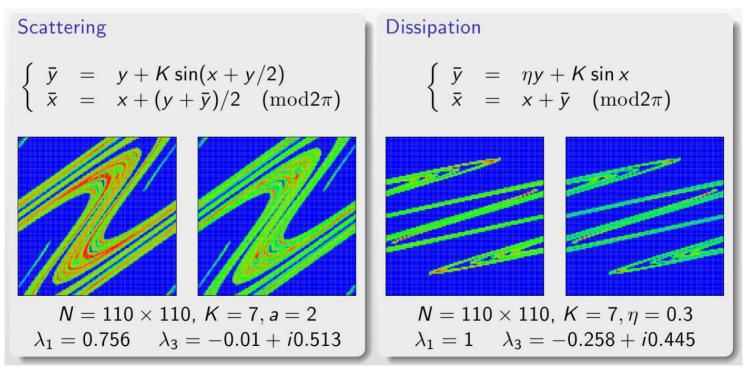
#### **Eigenvectors**





# Ulam method for dissipative systems

(Ermann, Shepelyansky (2010))



#### Fractal Weyl law

 $N_{\gamma} =$  number of Gamow eigenstates that have escape rates  $\gamma_j = -2 \ln |\lambda_j|$  in a finite bandwidth  $0 \le \gamma_j \le \gamma_b$ .

#### Fractal Weyl law for open quantum systems :

(e.g. Shepelyansky (2008))  $N_\gamma \propto N^{d-1} \propto \hbar^{-(d-1)}$  where d is a fractal dimension of a strange invariant set formed by orbits non-escaping in the future.

#### Fractal Weyl law for Ulam networks : $N_\gamma \propto N^ u \propto N^{d_0/2}$

 $d_0=$  dimension of invariant set of strange repeller (formed by orbits non-escaping in the future **and** in the past).  $\nu=d_0/2$ 

 $d = \operatorname{dimension} \operatorname{of} \operatorname{orbits} \operatorname{non-escaping} \operatorname{in}$  the future

$$d = d_0/2 + 1$$
 (inset)

## **University networks**

(KF, Georgeot, Shepelyansky (2011))

In realistic WWW networks invariant subspaces of nodes create large degeneracies of  $\lambda_1$  (or  $\lambda_2$  if  $\alpha < 1$ ) which is very problematic for the Arnoldi method.

Therefore determine the *invariant subspaces* as follows:

Let  $N_c = bN$  a certain fraction of the network size N (e.g. b = 0.1).

- For a given initial node  $i_0$  determine a sequence of node sets  $S_n$  by  $S_0 = \{i_0\}$  and  $S_{n+1}$  is the set containing all nodes of  $S_n$  and those which can be reached by a link from a node in  $S_n$ .
- If  $S_n = S_{n+1}$  with at most  $N_c$  elements for some  $n \Rightarrow S_n$  is an *invariant subspace*.

- If for some n the set  $S_n$  contains a dangling node (connected by construction to any other node) or if  $S_n$  contains more than  $N_c$  elements  $\Rightarrow i_0$  is identified as a node belonging to the *core space* (space of nodes not belonging to an invariant subspace).
- Repeat the procedure for every network node as potential initial node except for those nodes which are already identified as subspace nodes. If for some *n* the set  $S_n$  contains a previously found core space node  $\Rightarrow i_0$  also belongs to the core space.
- Merge all subspaces with common members. In this way one obtains a decomposition of the network in many *separate subspaces* with  $N_s$  nodes and a "big" *core space*.

This procedure can be efficiently implemented as a computer program. It turns out that for most networks the exact choice of b is not important (e.g. b = 0.1 or b = 0.9) as long as b = O(1). Note that a core space node may have a link to an invariant subspace but a subspace node may not have a link to another subspace or the core space.

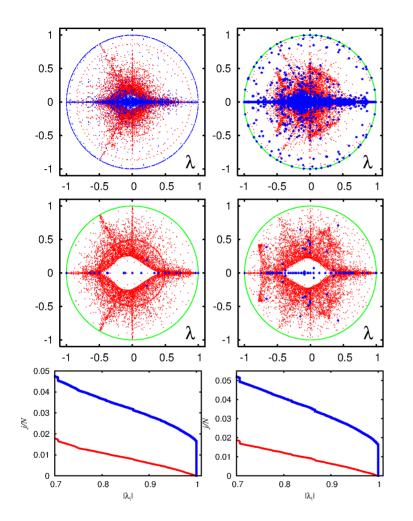
The decomposition in subspaces and a core space implies a block structure of the matrix S:

$$S = \left(\begin{array}{cc} S_{ss} & S_{sc} \\ 0 & S_{cc} \end{array}\right)$$

where  $S_{ss}$  is block diagonal according to the subspaces. The subspace blocks of  $S_{ss}$  are all matrices of PF type with at least one eigenvalue  $\lambda_1 = 1$  explaining the high degeneracies.

To determine the spectrum of S apply:

- Exact (or Arnoldi) diagonalization on each subspace.
- The Arnoldi method to  $S_{cc}$  to determine the largest core space eigenvalues  $\lambda_j$  (note:  $|\lambda_j| < 1$ ). The largest eigenvalues of  $S_{cc}$ are no longer degenerate but other degeneracies are possible (e.g.  $\lambda_j = 0.9$  for Wikipedia).



Cambridge 2006 (left),  $N = 212710, N_s = 48239$ 

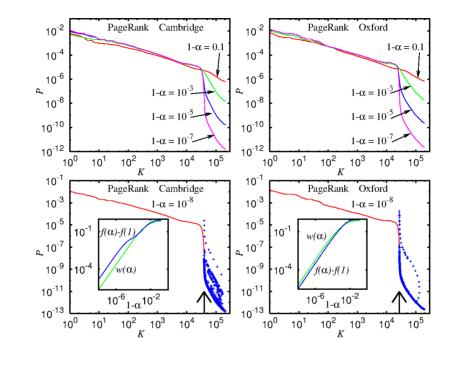
Oxford 2006 (right), N = 200823,  $N_s = 30579$ 

Spectrum of S (upper panels),  $S^*$  (middle panels) and dependence of rescaled level number on  $|\lambda_j|$  (lower panels).

Blue: subspace eigenvalues Red: core space eigenvalues (with Arnoldi dimension  $n_A = 20000$ )

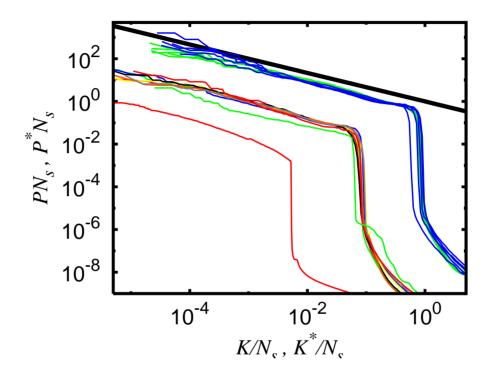
University networks

#### PageRank for $\alpha \to 1$ :



$$P = \sum_{\substack{\lambda_j=1 \\ \text{subspace contributions}}} c_j \psi_j + \sum_{\substack{\lambda_j\neq 1 \\ \lambda_j\neq 1}} \frac{1-\alpha}{(1-\alpha) + \alpha(1-\lambda_j)} c_j \psi_j .$$

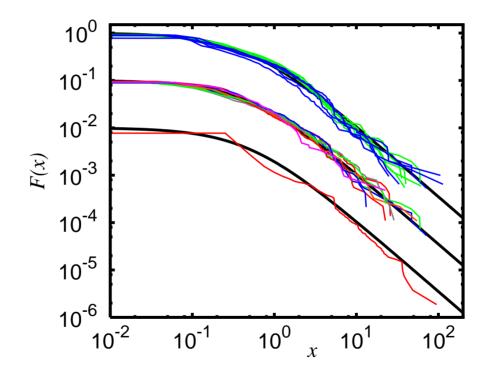
Rescaled PageRank at  $\alpha = 1 - 10^{-8}$  :



Top: Cambridge, Oxford 2002-2006; middle: other universities; bottom: Wikipedia<sup>\*</sup>; black line  $\propto K^{-2/3}$ ;  $N_s =$  sum of all subspace dimensions.

#### **Distribution of dimensions of invariant subspaces**

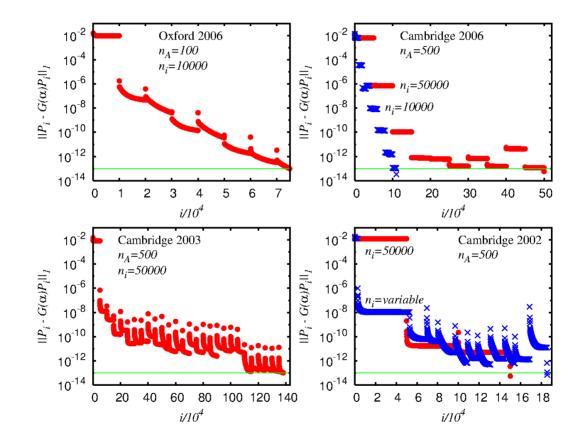
F(x) = fraction of invariant subspaces with dimension larger than  $x\langle d \rangle$  where  $\langle d \rangle =$  average subspace dimension.



Top: Cambridge, Oxford 2002-2006; middle: other universities; bottom: Wikipedia\*; black line:  $F(x) = 1/(1+2x)^{3/2}$ .

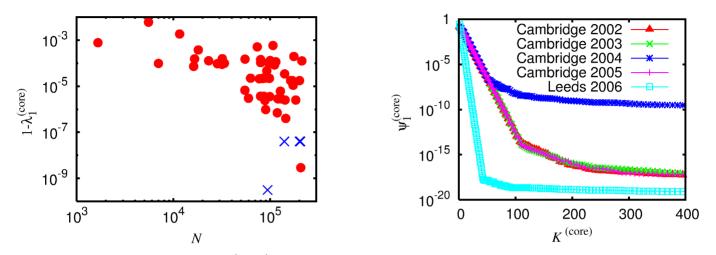
#### Numerical PageRank method for $\alpha \to 1$

Combination of power method and Arnoldi diagonalization :



Here:  $\alpha = 1 - 10^{-8}$ 





Left: Core space gap  $1 - \lambda_1^{(\text{core})}$  vs N for certain british universities. Red dots for gap  $> 10^{-9}$ ; blue crosses (moved up by  $10^9$ ) for gap  $< 10^{-16}$ .

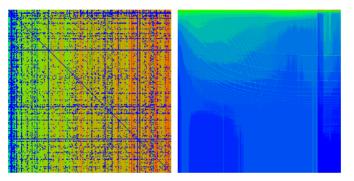
Right: first core space eigenvecteur for universities with gap  $< 10^{-16}$  or gap  $= 2.91 \times 10^{-9}$  for Cambridge 2004.

Core space gaps  $< 10^{-16}$  correspond to *quasi-subspaces* where it takes quite many "iterations" to reach a dangling node.

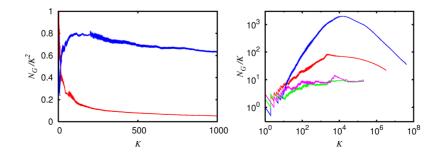
### **Twitter network**

(KF, Shepelyansky (2012), preprint) Twitter 2009 : N = 41652230 nodes,  $N_{\ell} = 1468365182$  network links.

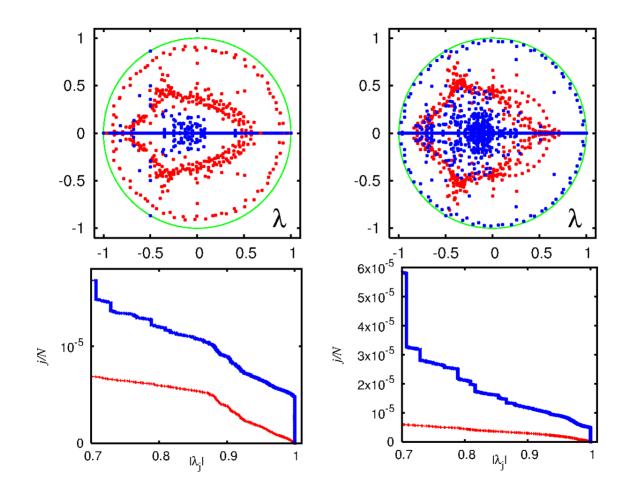
Matrix structure in K-rank order:



Number  $N_G$  of non-empty matrix elements in  $K \times K$ -square:

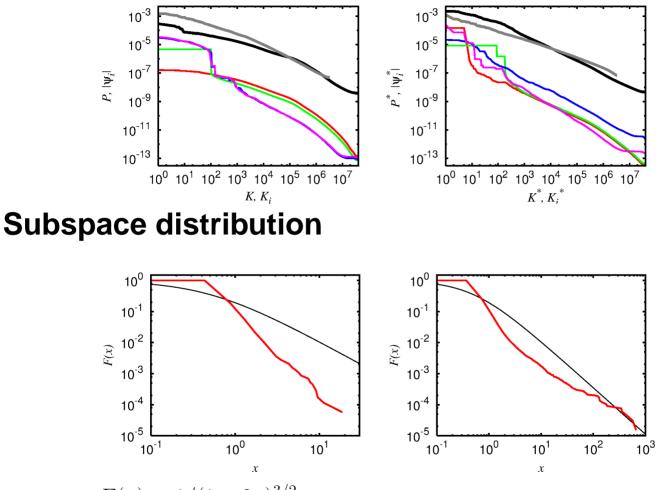


#### Spectrum



 $n_A = 640 \implies 250 \text{ GB of RAM memory.}$ 

#### PageRank, CheiRank, eigenvectors

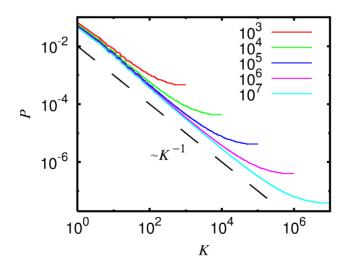


Black line:  $F(x) = 1/(1+2x)^{3/2}$ .

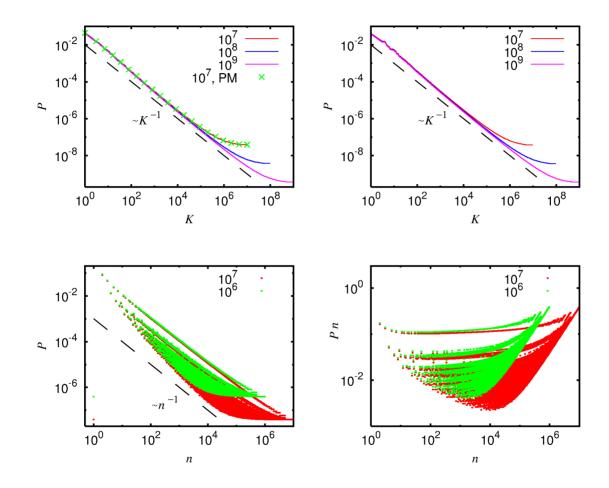
### Integer network

(KF, Chepelianskii, Shepelyansky (2012), preprint)

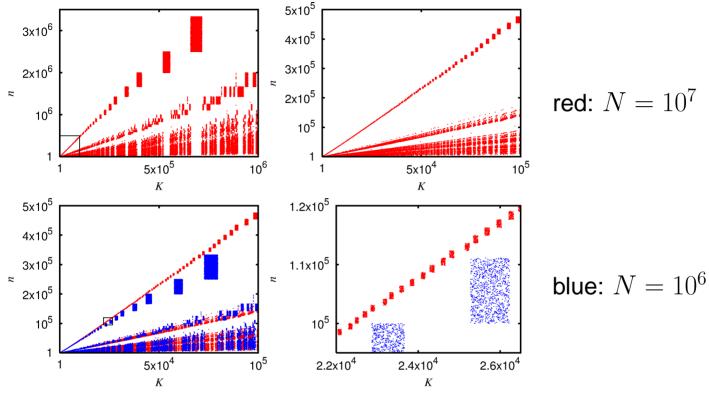
Consider the integers  $n \in \{1, ..., N\}$  and construct an adjacency matrix by  $A_{mn} = k$  where k is the largest integer such that  $m^k$  is a divisor of n if 1 < m < n and  $A_{mn} = 0$  if m = 1 or m = n (note  $A_{mn} = k = 0$  if m is not a divisor of n). Construct S and G in the usual way from A.



#### PageRank



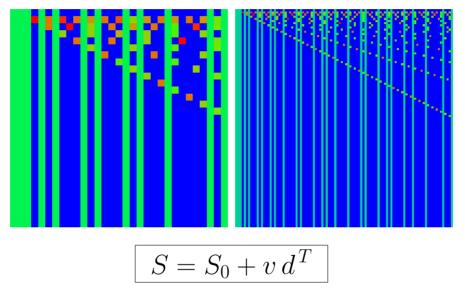
#### Dependence of $\boldsymbol{n}$ on K-index



"New order" of integers:  $K = 1, 2, ..., 32 \Rightarrow n = 2, 3, 5, 7, 4, 11, 13, 17, 6, 19, 9, 23, 29, 8, 31, 10, 37, 41, 43, 14, 47, 15, 53, 59, 61, 25, 67, 12, 71, 73, 22, 21.$ 

## Semi-analytical determination of spectrum, PageRank and eigenvectors

Matrix structure:



where v = e/N,  $d_j = 1$  for dangling nodes (primes and 1) and  $d_j = 0$  otherwise.  $S_0$  is the pure link matrix which is *nil-potent*:

$$S_0^l = 0$$

with  $l = [\log_2(N)] \ll N$ .

Let  $\psi$  be an eigenvector of S with eigenvalue  $\lambda$  and  $C = d^T \psi$ .

- If  $C = 0 \Rightarrow \psi$  eigenvector of  $S_0 \Rightarrow \lambda = 0$  since  $S_0$  nil-potent.
- If  $C \neq 0 \Rightarrow \lambda \neq 0$  since the equation  $S_0 \psi = -C v$  does not have a solution  $\Rightarrow \lambda \mathbf{1} S_0$  invertible.

$$\Rightarrow \psi = C \left(\lambda \mathbf{1} - S_0\right)^{-1} v = \frac{C}{\lambda} \sum_{j=0}^{l-1} \left(\frac{S_0}{\lambda}\right)^j v$$

From 
$$\lambda^l = (d^T \psi/C) \lambda^l \Rightarrow \mathcal{P}_r(\lambda) = 0$$

with the *reduced polynomial* of degree  $l = [\log_2(N)]$ :

$$\mathcal{P}_r(\lambda) = \lambda^l - \sum_{j=0}^{l-1} \lambda^{l-1-j} c_j = 0 \quad , \quad c_j = d^T S_0^j v \; .$$

 $\Rightarrow$  at most l eigenvalues  $\lambda \neq 0$  which can be numerically determined as the zeros of  $\mathcal{P}_r(\lambda)$ . (Note:  $l \leq 29$  for  $N \leq 10^9$ ).

Furthermore for  $\lambda = 1 \Rightarrow$  PageRank:

$$P = C \sum_{j=0}^{l-1} S_0^j v , \ C = d^T P .$$

The subspace of  $\lambda \neq 0$  is represented by the vectors  $v^{(j)} = S_0^j v$  for  $j = 0, \ldots, l-1$ 

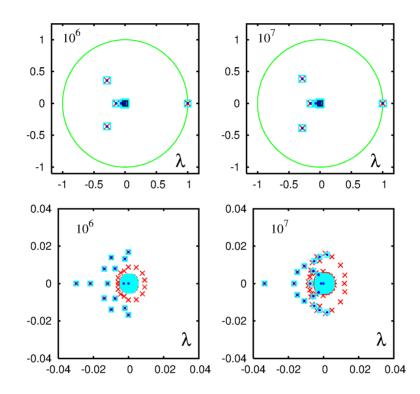
$$\Rightarrow \quad S \, v^{(j)} = c_j \, v^{(0)} + v^{(j+1)} = \sum_{k=0}^{l-1} \bar{S}_{k+1,j+1} \, v^{(k)}$$

"Small"  $l \times l$ -representation matrix :

$$\bar{S} = \begin{pmatrix} c_0 & c_1 & \cdots & c_{l-2} & c_{l-1} \\ 1 & 0 & \cdots & 0 & 0 \\ 0 & 1 & \cdots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & 1 & 0 \end{pmatrix} , \quad \bar{P} = C \begin{pmatrix} 1 \\ 1 \\ 1 \\ \vdots \\ 1 \end{pmatrix}$$
$$P = \sum_j \bar{P}_{j+1} v^{(j)} = C \sum_j v^{(j)} \text{ and due to sum rule: } \sum_j c_j = 1.$$

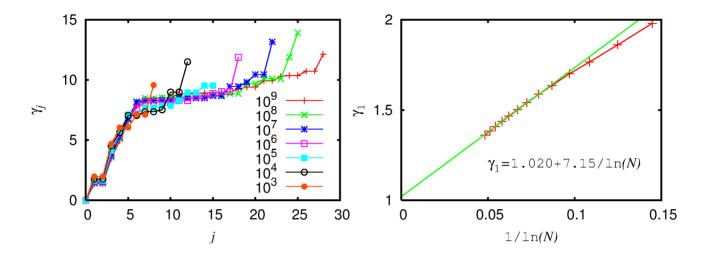
with

#### **Spectrum I**



blue dots: semi-analytical eigenvalues as zeros from  $\mathcal{P}_r(\lambda)$  (or eigenvalues of  $\overline{S}$ ). red crosses: Arnoldi method with random initial vector and  $n_A = 1000$ . light blue boxes: Arnoldi method with constant initial vector v = e/N and  $n_A = 1000$ .

#### **Spectrum II**



 $\gamma_j = -2\ln|\lambda_j|$ 

Large N limit of  $\gamma_1$  with the scaling parameter:  $1/\ln(N).$  Note:

$$c_0 = d^T v = \frac{1}{N} \sum_{j=1}^N d_j = \frac{1 + \pi(N)}{N} \approx \frac{1}{\ln(N)}$$

where  $\pi(N)$  is the number of primes below N.

## References

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